The Halting Problem Lecture 33 Section 12.1

Robb T. Koether

Hampden-Sydney College

Mon, Nov 14, 2016

Robb T. Koether (Hampden-Sydney College)

The Halting Problem

Mon, Nov 14, 2016 1 / 18

э

DQC

イロト イヨト イヨト イヨト









프 🖌 🖌 프

Image: A matrix

Outline

Decision Problems

2 The Halting Problem

3 Reduction of Problems

Assignment

э

DQC

イロト イヨト イヨト イヨト

An algorithm is a Turing machine that halts on every possible input.

э

イロト イポト イヨト イヨト

A decision problem is a problem that is clearly stated and unambiguously has a yes-or-no answer.

< ロト < 同ト < ヨト < ヨト

A decision problem is decidable if there is an algorithm that decides it. Otherwise, the problem is undecidable. That is, a decision problem is decidable if there is a Turing machine that can give the correct yes-or-no answer for every instance of the problem.

イロト イポト イヨト イヨト



2 The Halting Problem

3 Reduction of Problems

Assignment

Robb T. Koether (Hampden-Sydney College)

э

DQC

<ロト < 回ト < 回ト < 回ト

Definition (The Halting Problem)

The Halting Problem asks whether there exists a Turing machine H which, when fed as input a representation of a Turing machine M and an input w, will determine whether M would halt if it were executed with input w.

< ロト < 同ト < ヨト < ヨト

Theorem (The Halting Problem)

The Halting Problem is undecidable.

Robb T. Koether (Hampden-Sydney College)

э

イロト イポト イヨト イヨト

Proof.

- Suppose that the Halting Problem is decidable.
- Then there exists a Turing machine *H* that will decide the Halting Problem.
- We will build a Turing machine *M'* that does the following.
 - Read a representation of a Turing machine *M*.
 - Feed *M* and a copy of *M* as both machine and input to *H* to decide whether *M* would halt if it were executed with *M* as its own input.
 - If H reports "no," then M' reports "yes."
 - If *H* reports "yes," then *M*' loops.

A B F A B F

Proof.

- Now run *M*' and feed it a copy of itself as both the Turing machine and its input.
- M' must halt or loop.
- If M' reports "yes," then H must have reported "no," meaning that M' does not halt.
- But *M'* did halt, so that is a contradiction.
- If M' loops, then it is because H reported "yes," meaning that M' halts.
- But *M*′ did not halt, so that, too, is a contradiction.
- The conclusion is that no such Turing maching *H* exists.

∃ ► < ∃ ►</p>

Outline







4 Assignment

3

DQC

イロト イヨト イヨト イヨト

Definition (Reduction of Problems)

Given two decision problems A and B, a reduction of A to B is an algorithm (Turing machine) that will convert any instance of A into an instance of B such that the two instances have the same yes-or-no answer.

Given a Turing machine *M* and input *w*, the membership problem (or acceptance problem) for Turing machines is the question of whether $w \in L(M)$ (whether *M* accepts *w*).

Given a Turing machine *M* and input *w*, the membership problem (or acceptance problem) for Turing machines is the question of whether $w \in L(M)$ (whether *M* accepts *w*).

• We can reduce the membership problem to the halting problem.

Given a Turing machine *M* and input *w*, the membership problem (or acceptance problem) for Turing machines is the question of whether $w \in L(M)$ (whether *M* accepts *w*).

- We can reduce the membership problem to the halting problem.
- Can we reduce the halting problem to the membership problem?

- Let *A* be a decision problem and let *H* be the halting problem.
- If *H* can be reduced to *A*, then *A* is undecidable.

∃ ► < ∃ ►</p>

Given a Turing machine M, input w, and a state q of M, the state-entry problem asks whether M ever enters state q when processing w.

< ロト < 同ト < ヨト < ヨト

Given a Turing machine M, input w, and a state q of M, the state-entry problem asks whether M ever enters state q when processing w.

• The halting problem can be reduced to the state-entry problem.

Given a Turing machine M, input w, and a state q of M, the state-entry problem asks whether M ever enters state q when processing w.

- The halting problem can be reduced to the state-entry problem.
- Therefore, the state-entry problem is undecidable.

Outline

Decision Problems

- 2 The Halting Problem
- 3 Reduction of Problems



3

DQC

<ロト < 回ト < 回ト < 回ト

Homework

• Section 12.1 Exercises 2, 3, 7, 8, 12, 13.

Robb T. Koether (Hampden-Sydney College)

3

DQC

イロト イヨト イヨト イヨト